	Introduction
UNIX Network Programming	<ul> <li>Standard RPC mechanisms offer point-to- point semantics</li> </ul>
An Overview of IP Multicasting	<ul> <li>Many applications require more flexible com- munication semantics</li> </ul>
Douglas C. Schmidt	<ul> <li><i>e.g.</i>, 1-to-n or n-to-m communication</li> <li>Group communication is an abstraction that supports more flexible communication</li> </ul>
	<ul> <li>May be supported in hardware and/or software</li> </ul>
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Design Issues for Group	
Communication	Addressing
• Addressing	• Four methods of addressing:
• Reliability	<ol> <li>Sender explicitly enumerates addresses</li> <li>Single group address</li> </ol>
• Ordering	3. Source addressing
• Delivery semantics	4. Functional addressing
• Response semantics	<ul> <li>The latter two mechanisms are difficult to implement in hardware</li> </ul>
• Group Structure	
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### Reliability

- Reliability deals with recovery from communication failures
  - e.g., buffer overflows and bit errors
- Much more difficult to implement reliability for group communication efficiently
- Thus, many group communication mechanisms are unreliable

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#### Ordering

- Four ordering semantics:
- 1. No ordering
  - Easy to implement, but hard to use
- 2. FIFO ordering
  - Message from one member are delivered in order sent
- 3. Causal ordering
  - All messages that are related are ordered
- 4. Total ordering
  - Difficult to implement, but easy to use

#### **Delivery Semantics**

- Defines when a message is considered delivered successfully to a group
- Three choices
- 1. k-delivery
  - Succeeds when k participants have received message
- 2. Quorum delivery
  - $-\,$  Succeeds when a majority have received message
- 3. Atomic delivery
  - All must receive or none receive

#### **Response Semantics**

- Defines what a sender expects from the receivers
- Four choices
- 1. No responses
- 2 A single response
- 3. Many responses
- 4. All responses
- This is related to the reliability dimension...

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## Group Structure

- Two choices
- 1. Closed groups
  - Only group members can send to the group
- 2. Open groups
  - Non-members can also send to the group
- Two more choices
- 1. Static groups
  - Members must join for the duration of the group
- 2. Dynamic groups
  - Members can join and leave as they wish
- Note that support for *overlapping* groups is also available

### **IP** Multicasting

- IP multicasting implements a simple model of group communication
- IP multicasting is used to send an IP or UDP datagram to a finite number of hosts using a single IP address
- IP multicasting is becoming available on modern operating systems
  - e.g., Solaris and Windows NT

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# Advantages of IP Multicasting over Broadcasting

- Provides the ability to transmit information to a group of hosts via a single multicast address
- Multicasting is more efficient than broadcasting
  - $-\,$  It doesn't disturb hosts that are not participating in the communication
  - $\it i.e.,$  it reduces extraneous packet examination in lower-level protocols...

## **IP** Multicast Model

- Addressing
  - The IP address used for multicasting identifies a "multicast group"
    - ▷ A group may consist of 0 or more hosts
- Reliability, Delivery, and Reponse Semantics
  - IP multicasting is unreliable, delivery is on a best effort basis
- Ordering
  - IP multicasting provides no ordering guarantees

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IP Multicast Model (cont'd)	<ul> <li>IP Multicast Groups</li> <li>Two types of multicast groups</li> </ul>
<ul> <li>Group Structure</li> <li>A single host may belong to many groups</li> <li>A host may join or leave a group at any time during the lifetime of the group</li> <li>A host need not be a member of a group to send a message to the group</li> <li>Groups may overlap</li> </ul>	<ol> <li>Permanent         <ul> <li>Can have zero or more hosts with well-known IP address (e.g., 224.0.0.1)</li> <li>Transient                 <ul></ul></li></ul></li></ol>
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<ul> <li>IP Multicast Addressing</li> <li>Accomplished through the use of IP class D addresses</li> <li>Four high order bits are set to 1110</li> <li>Remaining 28 bits identify specific multicast groups</li> <li>▶ Ranges from 224.0.0.0 to 239.255.255.255</li> <li>Address 224.0.0.0 is permanently unassigned</li> <li>Address 224.0.0.1 is a permanent group to address all multicast hosts on a directly connected subnetwork</li> </ul>	<ul> <li>Mapping Multicast Addresses to Ethernet Addresses</li> <li>Multicast IP packets ultimately resolve down to Ethernet destinations</li> <li>To create a unique Ethernet address the low or- der 23 bits of the IP address are mapped to the Ethernet address</li> <li>Note the potential for conflict since there are 28 significant bits in a Class D IP address, but only the lower 23 are used</li> </ul>
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# **IP** Multicast Administration (IGMP) **Host Participation** • The Internet Group Management Protocol (IGMP) handles administrative tasks related • A host can be in one of three states when to hosts and gateways participating on a multicasting network: - IGMP resides upon the IP layer 1 Participate fully by belonging to a multicast group and sending/receiving multicast datagrams - Similar in spirit to ICMP, *i.e.*, sent in IP datagrams 2. Be configured to send, but not receive, multicast datagrams • IGMP keeps hosts/gateways informed about status and configuration of multicast groups 3. Not participate in mulitcasting at all - Accomplished by continuously querying hosts and waiting for responses to be sent by one of the hosts in the group 17 18 **Programming IP Multicast via** Sockets **Example Header File** • Use setsockopt to join/leave a multicast group: • Shared by both sender and receiver struct ip\_mreq mcast\_addr; // Join group // Multicast group address. setsockopt (sd, IPPROTO\_IP, IP\_ADD\_MEMBERSHIP, #define MCAST\_ADDR "224.9.9.2" (char \*) &mcast\_addr, sizeof mcast\_addr); // Port number. // Leave group #define UDP\_PORT 3112 setsockopt (sd, IPPROTO\_IP, IP\_DROP\_MEMBERSHIP, (char \*) &mcast\_addr, sizeof mcast\_addr); // Name of the interface. #define INTERFACE "le0" • See the ip(7) manual page for more info on options 20 19

## Example Receiver

• Server program

int main (void) {
 struct sockaddr\_in local;
 struct ifreq interface\_addr;
 struct in\_addr in\_addr;
 struct ip\_mreq mcast\_addr;
 struct sockaddr\_in sa;
 int sd = socket (PF\_INET, SOCK\_DGRAM, 0);

// Determine interface address.
strcpy (interface\_addr.ifr\_name, INTERFACE);
ioctl (sd, SIOCGIADDR, &interface\_addr);

mcast\_addr.imr\_multiaddr.s\_addr = inet\_addr (MCAST\_ADDR); sa = &interface\_addr.ifr\_addr; mcast\_addr.imr\_interface.s\_addr = sa->sin\_addr.s\_addr;

// Initialize local port number. memset ((void \*) &local, 0, sizeof local); local.sin\_family = AF\_INET; local.sin\_port = htons (UDP\_PORT); local.sin\_addr.s\_addr = hton1 (INADDR\_ANY);

bind (sd, &local, sizeof local); // Go into loop on readfrom() and print what is received...

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#### **Example Sender**

• Multicast data to receiver(s)

int main (int argc, char \*argv[]) {
 struct sockaddr\_in sa;
 struct hostent \*hp;
 int sd = socket (AF\_INET, SOCK\_DGRAM, 0);
 sa.sin\_family = AF\_INET;
 sa.sin\_port = htons (UDP\_PORT);
 sa.sin\_addr.s\_addr = inet\_addr (MCAST\_ADDR);
 hp = gethostbyname (argv[1]);
 memcpy (&sa.sin\_addr, hp->h\_addr, hp->h\_length);

// Multicast data to receiver(s) via sendto();

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## **UNIX** Utilities

- netstat
- View multicast groups via netstat -g
- View multicast statistics via netstat -s
- snoop
- View live packets as they arrive via  ${\tt snoop}\ {\tt multicast}$
- See manual pages for options...

#### Summary

- Group communication is increasingly important
- IP multicast provides very low-level mechanisms to support group communication
- Building more sophisticated semantics on top of IP multicasting remains a research issue

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